

Height Functions and Extremal Matchings on General Planar Bipartite Graphs

LP

1 Setup

We work with **planar bipartite graphs** — connected bipartite graphs equipped with a fixed embedding in the plane (or the sphere). Height functions are defined on *faces* of the embedding, so planarity is essential: without it there is no well-defined facial structure (e.g., $K_{3,3}$ is not planar, and height functions make no sense there). The general theory goes back to Thurston [Thu90] for domino tilings and extends to arbitrary planar bipartite graphs; see Kenyon [Ken09], Propp [Pro03].

2 Superposition and alternating cycles

Fix a *reference perfect matching* M_{ref} . Given any other perfect matching M , the symmetric difference

$$M \oplus M_{\text{ref}} = (M \setminus M_{\text{ref}}) \cup (M_{\text{ref}} \setminus M) \quad (2.1)$$

has degree 0 or 2 at every vertex, so it decomposes into disjoint *alternating cycles* whose edges alternate between M and M_{ref} .

To extract a height function from this data, we impose a canonical orientation: direct every M -edge from **black** to **white**, and every M_{ref} -edge from **white** to **black**. Each alternating cycle then becomes a consistently directed closed curve in the plane.

We illustrate with the 2×4 grid graph G (three square faces, five perfect matchings M_1, \dots, M_5). Taking $M_{\text{ref}} = M_5$ (all vertical edges) and $M = M_2$ (all horizontal), the symmetric difference produces two disjoint clockwise alternating cycles (Figure 1).

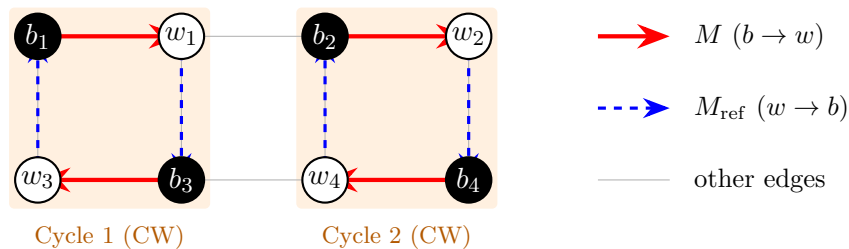


Figure 1: Superposition of $M = M_2$ (horizontal, red) and $M_{\text{ref}} = M_5$ (vertical, blue) on G . The symmetric difference produces two disjoint clockwise alternating cycles (shaded), one around each end face. The middle face’s horizontal boundary edges belong to neither matching.

3 The height function

Set $h = 0$ on the outer (unbounded) face. For two adjacent faces f, f' sharing a directed edge e on their boundary, set

$$h(f') = h(f) + 1 \quad \text{if } e \text{ flows to the left of the crossing path,} \quad (3.1)$$

and $h(f') = h(f) - 1$ if e flows to the right. Edges in $M \cap M_{\text{ref}}$ carry no directed flow, and the height does not change across them. Well-definedness follows from planarity — the consistent orientation of alternating cycles guarantees zero height gain around any closed walk.

Figure 2 shows the resulting height functions for two matchings of G , with faces colored by height value (green for positive, red for negative, gray for zero).



Figure 2: Height functions for two matchings of G with $M_{\text{ref}} = M_5$. Matching edges shown in red. The height vector (h_L, h_M, h_R) records the height of the left, middle, and right face.

4 Local flips

A face is *flippable* if its entire boundary forms an alternating cycle of M - and M_{ref} -edges (equivalently, a single directed cycle under the canonical orientation). In lozenge tilings this is a hexagonal face; in domino tilings, a 2×2 square. The cycle direction determines the effect of the flip:

- **Clockwise** boundary — local height maximum; flipping (swapping M -edges for M_{ref} -edges on the face) strictly *decreases* the height.
- **Counter-clockwise** boundary — local height minimum; flipping strictly *increases* the height.

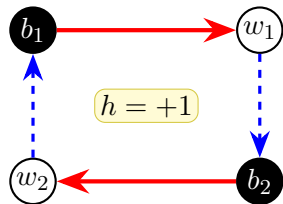


Figure 3: Before flip: clockwise cycle (local height maximum).

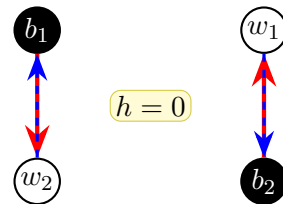


Figure 4: After flip: M and M_{ref} now agree on the vertical edges (both arrows on each edge), so the alternating cycle disappears and h drops from $+1$ to 0 .

After the flip, M and M_{ref} agree on the flipped face, so the alternating cycle there disappears. The net effect is that the height of the flipped face changes by ± 1 .

5 Extremal matchings and the distributive lattice

Starting from any perfect matching M , we can construct the unique *maximal matching* M_{\max} by repeatedly flipping all faces with clockwise alternating boundaries (local height minima), and the unique *minimal matching* M_{\min} by flipping all counter-clockwise boundaries (local height maxima).

The flip order does not matter. The set of all perfect matchings, partially ordered by pointwise comparison of height functions, forms a **distributive lattice** — see Thurston [Thu90], Propp [Pro03], and the general theory of \mathbb{Z} -flows on planar graphs. In particular:

- The lattice has a unique maximum M_{\max} and a unique minimum M_{\min} .
- Every increasing chain of flips terminates at M_{\max} , regardless of order; every decreasing chain terminates at M_{\min} .
- The lattice operations $M_1 \vee M_2$ and $M_1 \wedge M_2$ correspond to pointwise max and min of the height functions.

This confluent structure — no local traps, no dependence on flip sequence — is the key reason height functions are such a powerful tool for the analysis of the space of perfect matchings of planar bipartite graphs.

For the graph G , the five perfect matchings and their height vectors (h_L, h_M, h_R) are shown in Figure 5.

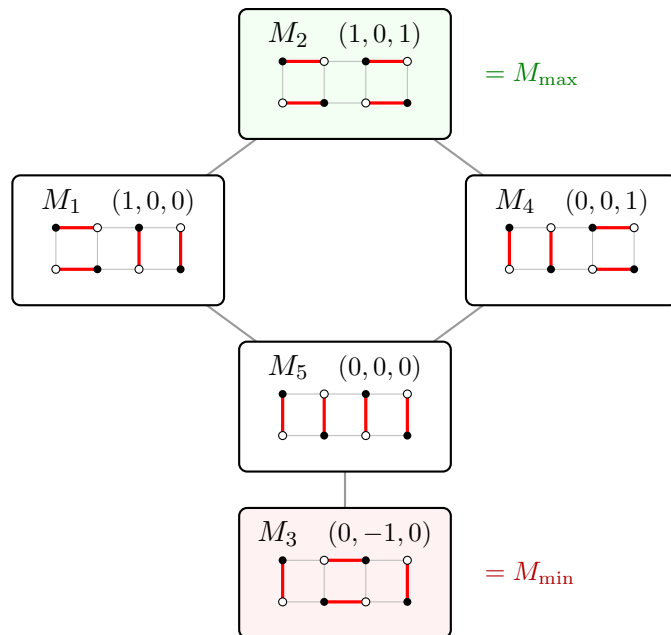


Figure 5: The distributive lattice of all five perfect matchings of G , partially ordered by pointwise comparison of height functions. Each node shows the matching name, height vector (h_L, h_M, h_R) , and a thumbnail (red edges = matching). Edges of the Hasse diagram connect matchings that differ by a single face flip.

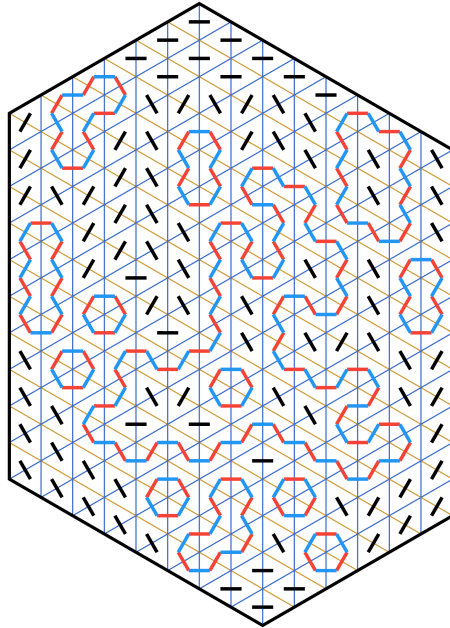


Figure 6: A uniform random double-dimer configuration on a large lozenge graph ($n = 188$), illustrating the macroscopic loop structure that emerges from the superposition of two independent perfect matchings.

References

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- [Pro03] J. Propp. Generalized domino-shuffling. *Theoretical Computer Science*, 303(2-3):267–301, 2003. arXiv:math/0111034 [math.CO]. 1, 3
- [Thu90] W. P. Thurston. Conway’s tiling groups. *Amer. Math. Monthly*, 97:757–773, 1990. 1, 3